

Note

## Kernels in quasi-transitive digraphs

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### Abstract

Let  $D$  be a digraph,  $V(D)$  and  $A(D)$  will denote the sets of vertices and arcs of  $D$ , respectively.

A kernel  $N$  of  $D$  is an independent set of vertices such that for every  $w \in V(D) - N$  there exists an arc from  $w$  to  $N$ . A digraph is called *quasi-transitive* when  $(u, v) \in A(D)$  and  $(v, w) \in A(D)$  implies  $(u, w) \in A(D)$  or  $(w, u) \in A(D)$ . This concept was introduced by Ghoulilá-Houri [Caractérisation des graphes non orientés dont on peut orienter les arrêtes de maniere à obtenir le graphe d' un relation d'ordre, C.R. Acad. Sci. Paris 254 (1962) 1370–1371] and has been studied by several authors. In this paper the following result is proved: Let  $D$  be a digraph. Suppose  $D = D_1 \cup D_2$  where  $D_i$  is a quasi-transitive digraph which contains no asymmetrical infinite outward path (in  $D_i$ ) for  $i \in \{1, 2\}$ ; and that every directed cycle of length 3 contained in  $D$  has at least two symmetrical arcs, then  $D$  has a kernel. All the conditions for the theorem are tight.

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### 1. Introduction

For general concepts we refer the reader to [4]. In the paper we write digraph to mean 1-digraph in the sense of Berge [4]. In this paper  $D$  will denote a possibly infinite digraph with  $V(D)$  and  $A(D)$  being the sets of vertices and arcs of  $D$ , respectively. Often we shall write  $u_1u_2$  instead of  $(u_1, u_2)$ . An arc  $u_1u_2 \in A(D)$  is called asymmetrical (resp. symmetrical) if  $u_2u_1 \notin A(D)$  (resp.  $u_2u_1 \in A(D)$ ). If  $S$  is a nonempty subset of  $V(D)$  then the subdigraph  $D[S]$  induced by  $S$  is the digraph with vertex set  $S$  and whose arcs are those arcs of  $D$  which join vertices of  $S$ .

A directed path is a finite or infinite sequence  $(x_1, x_2, \dots)$  of distinct vertices of  $D$  such that  $(x_i, x_{i+1}) \in A(D)$  for each  $i$ . When  $D$  is infinite and the sequence is infinite we call the directed path an infinite outward path. Let  $S_1$  and  $S_2$  be subsets of  $V(D)$ . A finite directed path  $(x_1, \dots, x_n)$  will be called an  $S_1S_2$ -directed path whenever  $x_1 \in S_1$  and  $x_2 \in S_2$ , in particular when the directed path is an arc, we will call it an  $S_1S_2$ -arc.

**Definition 1.1.** A set  $I \subseteq V(D)$  is independent if  $A(D[I]) = \emptyset$ . A kernel  $N$  of  $D$  is an independent set of vertices such that for each  $z \in V(D) - N$  there exists a  $zN$ -arc in  $D$ .

A digraph  $D$  is called a kernel-perfect digraph when every induced subdigraph of  $D$  has a kernel.

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The concept of kernel was introduced by Von Neumann and Morgenstern [15] in the context of Game Theory. The problem of the existence of a kernel in a given digraph has been studied by several authors in particular by Richardson [16,17], Duchet and Meyniel [9], Duchet [7,8], Galeana-Sánchez and Neumann-Lara [10].

A digraph  $D$  is transitive whenever  $(u, v) \in A(D)$  and  $(v, w) \in A(D)$  implies  $(u, w) \in A(D)$ . A digraph is called *quasi-transitive* if whenever  $(u, v) \in A(D)$  and  $(v, w) \in A(D)$ , then  $(u, w) \in A(D)$  or  $(w, u) \in A(D)$ .

Quasi-transitive digraphs were introduced by Ghouilá-Houri [12] and have been studied by several authors for example Bang-Jensen and Huang [1–3], Huang [13], Skrien [19]. It was proved by Ghouilá-Houri [12] that an undirected graph can be oriented as a quasi-transitive digraph if and only if it can be oriented as a transitive digraph, namely a comparability graph. More information about comparability graphs can be found in [11,14].

In [6] Boros and Gurvich proved that if  $G$  is a perfect graph then any orientation of  $G$  in which each complete subdigraph has a kernel is kernel-perfect. It is well known that comparability graphs are perfect graphs (see for example [5]). Meyniel [9] observed that if  $D$  is a digraph such that every directed cycle of length 3 has at least two symmetrical arcs, then each complete subdigraph of  $D$  has a kernel.

We can conclude the following result.

**Theorem 1.2.** *If  $D$  is a finite quasi-transitive digraph such that every directed cycle of length 3 has at least two symmetrical arcs, then  $D$  is a kernel-perfect digraph.*

The result proved in this paper generalizes Theorem 1.2 and the following result of Sands et al. [18].

**Theorem 1.3** (Sands et al. [18]). *Let  $D$  be a digraph whose arcs are colored with two colors. If  $D$  contains no monochromatic infinite outward path, then there exists a set  $S$  of vertices of  $D$  such that no two vertices of  $S$  are connected by a monochromatic directed path and for every vertex  $x$  not in  $S$  there is a monochromatic directed path from  $x$  to a vertex in  $S$ .*

We include the following definitions in order to understand Theorem 1.3 in terms of kernels.

We call the digraph  $D$  an  $m$ -colored digraph if the arcs of  $D$  are colored with  $m$  colors. A directed path is called monochromatic if all of its arcs are colored alike. A kernel by monochromatic paths in an  $m$ -colored digraph  $D$  is a set of vertices  $N$  which satisfies the following two conditions: (i) for every pair of different vertices  $u, v \in N$  there is no monochromatic directed path between them; and (ii) for every vertex  $x \in V(D) - N$  there is a vertex  $y \in N$  such that there is an  $xy$ -monochromatic directed path.

If  $D$  is an  $m$ -colored digraph then the closure of  $D$ , denoted  $\mathcal{C}(D)$  is the digraph defined as follows:  $V(\mathcal{C}(D)) = V(D)$  and  $(u, v) \in A(\mathcal{C}(D))$  iff there exists a  $uv$ -monochromatic directed path contained in  $D$ .

Note that for any  $m$ -colored  $D$ ,  $D$  has a kernel by monochromatic paths if and only if  $\mathcal{C}(D)$  has a kernel.

In this terminology Theorem 1.3 asserts that if  $D$  is a 2-colored digraph, which contains no monochromatic infinite outward path, then  $\mathcal{C}(D)$  has a kernel.

Now it is clear that Theorem 1.3 is equivalent to the following assertion. Let  $D$  be a digraph;  $D_1$  and  $D_2$  transitive subdigraphs of  $D$  such that  $D = D_1 \cup D_2$  (recall that  $D_1 \cup D_2$  is defined as follows:  $V(D_1 \cup D_2) = V(D_1) \cup V(D_2)$  and  $A(D_1 \cup D_2) = A(D_1) \cup A(D_2)$ ) and  $A(D_1) \cap A(D_2) = \emptyset$ . If  $D$  has no infinite outward path contained in  $D_i$  ( $i = 1, 2$ ), then  $D$  has a kernel.

Finally, we will introduce some notation. Two subdigraphs  $D_1$  and  $D_2$  of  $D$  are given (possibly  $A(D_1) \cap A(D_2) \neq \emptyset$ ). For distinct vertices  $x, y$  of  $D$ ,  $x \xrightarrow{i} y$  will mean that the arc  $(x, y) \in A(D_i)$  and  $x \xrightarrow{i} S$  will mean that there exists an arc in  $D_i$  from  $x$  to a vertex in  $S$ , the negation of  $x \xrightarrow{i} y$  (resp.  $x \xrightarrow{i} S$ ) will be denoted by  $x \not\xrightarrow{i} y$  (resp.  $x \not\xrightarrow{i} S$ ), for  $i = 1, 2$ . When we do not know if the arc is in  $D_1$  or in  $D_2$  we write simply  $x \rightarrow y$ ; and  $x \rightarrow y$  will mean that  $(x, y) \notin A(D)$ . A directed cycle of length 3 will be called a triangle.

## 2. Kernels in the union of two quasi-transitive digraphs

The main result of this section is Theorem 2.3. The proof is similar to that in Sands et al. [18].

**Lemma 2.1.** *Let  $D$  be a digraph such that every triangle has at least two symmetrical arcs. If  $D_1$  is a quasi-transitive subdigraph of  $D$  and  $(v_1, v_2, \dots, v_n)$  is a sequence of vertices of  $D_1$  such that  $(v_i, v_{i+1}) \in A(D_1)$  and*

$(v_{i+1}, v_i) \notin A(D)$ , then the sequence is an asymmetrical directed path of  $D$  contained in  $D_1$ , and for each  $i \in \{1, \dots, n-1\}$ ,  $(v_i, v_{i+1}) \in A(D_1)$  and  $(v_j, v_i) \notin A(D)$  for every  $j \in \{i+1, \dots, n\}$ .

**Proof.** We proceed by induction on  $n$ . The result is obvious for  $n \leq 2$ . Assume the result is true for a sequence  $(v_1, \dots, v_n)$  which satisfies the hypothesis of Lemma 2.1. Consider a sequence  $T = (v_1, \dots, v_n, v_{n+1})$  such that for each  $i \in \{1, \dots, n\}$ ,  $(v_i, v_{i+1}) \in A(D_1)$  and  $(v_{i+1}, v_i) \notin A(D)$ . Since  $T' = (v_1, \dots, v_n)$  satisfies the inductive hypothesis, we have that  $T'$  is an asymmetrical directed path contained in  $D_1$  and for each  $i \in \{1, \dots, n-1\}$   $(v_i, v_j) \in A(D_1)$  and  $(v_j, v_i) \notin A(D)$  for every  $j \in \{i+1, \dots, n\}$ . So we only need to prove that for each  $i \in \{1, \dots, n-1\}$ ,  $v_i \neq v_{n+1}$ ,  $(v_i, v_{n+1}) \in A(D_1)$  and  $(v_{n+1}, v_i) \notin A(D)$ .

First assume by contradiction that  $v_{n+1} = v_i$  for some  $i \in \{1, \dots, n-1\}$ . It follows from the inductive hypothesis on  $T'$  that  $(v_i, v_n) = (v_{n+1}, v_n) \in A(D_1)$  and thus  $(v_{n+1}, v_n) \in A(D)$  contradicting our hypothesis on  $T$ . We conclude that  $T$  is an asymmetrical directed path of  $D$  contained in  $D_1$ . Now, we have from the inductive hypothesis on  $T'$  that for each  $i \in \{1, \dots, n-1\}$ ,  $(v_i, v_n) \in A(D_1)$  and since  $(v_n, v_{n+1}) \in A(D_1)$  and  $D_1$  is a quasi-transitive digraph, we have that  $(v_i, v_{n+1}) \in A(D_1)$  or  $(v_{n+1}, v_i) \in A(D_1)$ . If  $(v_{n+1}, v_i) \in A(D)$  then  $C_3 = (v_i, v_n, v_{n+1}, v_i)$  is a triangle and from the hypothesis on  $D$ ,  $C_3$  has at least two symmetrical arcs which is impossible as  $(v_{n+1}, v_n) \notin A(D)$  (hypothesis on  $T$ ) and  $(v_n, v_i) \notin A(D)$  (inductive hypothesis). Thus  $(v_{n+1}, v_i) \notin A(D)$  and  $(v_i, v_{n+1}) \in A(D_1)$ .  $\square$

**Lemma 2.2.** Let  $D$  be a digraph such that every triangle has at least two symmetrical arcs, and  $D_1$  be a quasi-transitive subdigraph of  $D$  which contains no asymmetrical (in  $D$ ) infinite outward path. If  $\emptyset \neq U \subseteq V(D)$  then there exists  $x \in U$  such that for all  $y \in U(x, y) \in A(D_1)$  implies  $(y, x) \in A(D)$ .

**Proof.** Suppose by contradiction that for each  $x \in U$ , there exists  $y \in U$  such that  $(x, y) \in A(D_1)$  and  $(y, x) \notin A(D)$ . Consider some  $x_1 \in U$ . Then there exists  $x_2 \in U$  such that  $(x_1, x_2) \in A(D_1)$  and  $(x_2, x_1) \notin A(D)$ . So for each  $n \in \mathbb{N}$ , given  $x_n \in U$ , there exists  $x_{n+1} \in U$  such that  $(x_n, x_{n+1}) \in A(D_1)$  and  $(x_{n+1}, x_n) \notin A(D)$ . It follows from Lemma 2.1 that  $T_{n+1} = (x_1, x_2, \dots, x_{n+1})$  is an asymmetrical directed path of  $D$  contained in  $D_1$ . Consider the sequence  $T = (x_n)_{n \in \mathbb{N}}$ ; for each  $n \in \mathbb{N}$ ,  $(x_n, x_{n+1}) \in A(D_1)$ , and for  $n < m$  we have  $\{x_n, x_m\} \subseteq V(T_m)$  and since  $T_m$  is a directed path we obtain  $x_n \neq x_m$ ; hence  $T$  is an asymmetrical infinite outward path of  $D$  contained in  $D_1$ , a contradiction.  $\square$

**Theorem 2.3.** Let  $D$  be a digraph such that  $D = D_1 \cup D_2$  (possibly  $A(D_1) \cap A(D_2) \neq \emptyset$ ), where  $D_i$  is a quasi-transitive subdigraph of  $D$  which contains no asymmetrical (in  $D$ ) infinite outward path. If every triangle contained in  $D$  has at least two symmetrical arcs, then  $D$  is a kernel-perfect digraph.

**Proof.** It suffices to prove that  $D$  has a kernel, as any induced subdigraph of  $D$  satisfies the hypothesis of Theorem 2.3.

For independent sets  $S, T$  of  $D$ , we write  $S \leq T$  if and only if for each  $s \in S$  there exists  $t \in T$  such that either  $s = t$  or  $(s \xrightarrow{1} t$  and  $t \rightarrow s)$ . Note that in particular  $S \subseteq T$  implies  $S \leq T$ .

(1) The collection of all independent sets of vertices of  $D$  is partially ordered by  $\leq$ .

(1.1)  $\leq$  is reflexive.

This follows from the fact  $S \subseteq S$ .

(1.2)  $\leq$  is transitive.

Let  $S, T$  and  $R$  be independent sets of vertices of  $D$ , such that  $S \leq T$  and  $T \leq R$ , and let  $s \in S$ . Since  $S \leq T$  there exists  $t \in T$  such that either  $s = t$  or  $(s \xrightarrow{1} t$  and  $t \rightarrow s)$  and  $T \leq R$  implies that there exists  $r \in R$  such that either  $t = r$  or  $(t \xrightarrow{1} r$  and  $r \rightarrow t)$ . If  $s = t$  or  $t = r$ , then  $s = r$  or  $(s \xrightarrow{1} r$  and  $r \rightarrow s)$  with  $r \in R$ . So we can assume  $s \neq t, t \neq r, (s \xrightarrow{1} t$  and  $t \rightarrow s)$  and  $(t \xrightarrow{1} r$  and  $r \rightarrow t)$ . Since  $D_1$  is a quasi-transitive digraph it follows from Lemma 2.1 on the sequence  $(s, t, r)$  that  $(s \xrightarrow{1} r$  and  $r \rightarrow s)$ .

(1.3)  $\leq$  is antisymmetrical.

Let  $S$  and  $T$  be independent sets of vertices of  $D$  such that  $S \leq T$  and  $T \leq S$ , and let  $s \in S$ . Since  $S \leq T$  there exists  $t \in T$  such that either  $s = t$  or  $(s \xrightarrow{1} t$  and  $t \rightarrow s)$ . Suppose that  $s \neq t$ . The fact  $T \leq S$  implies that there exists  $s' \in S$  such that either  $t = s'$  or  $(t \xrightarrow{1} s'$  and  $s' \rightarrow t)$ . When  $t = s'$  we obtain  $s \xrightarrow{1} s'$  contradicting that  $S$  is an independent set; so  $t \neq s'$  and  $(t \xrightarrow{1} s'$  and  $s' \rightarrow t)$ . Now applying Lemma 2.1 on the sequence  $(s, t, s')$ , we have  $s \xrightarrow{1} s'$  contradicting that  $S$  is an independent set. We conclude  $s = t$  and consequently  $s \in T$  and  $S \subseteq T$ . Analogously it can be proved  $T \subseteq S$ .

Let  $\mathcal{F}$  be the family of all nonempty independent sets  $S$  of vertices of  $D$  such that  $S \xrightarrow{2} y$  implies  $y \rightarrow S$  for all vertices  $y$  of  $D$ .

(2)  $(\mathcal{F}, \leq)$  has maximal elements.

(2.1)  $\mathcal{F} \neq \emptyset$ .

Since  $D_2$  is a quasi-transitive digraph which contains no asymmetrical infinite outward path, it follows from Lemma 2.2 (taking  $U = V(D)$  and  $D_2$  instead of  $D_1$ ) that there exists a vertex  $x \in V(D)$  such that  $x \xrightarrow{2} y$  implies  $y \rightarrow x$ , for all vertices  $y$  of  $D$ , so  $\{x\} \in \mathcal{F}$ .

(2.2) Every chain in  $(\mathcal{F}, \leq)$  is upper bounded.

Let  $\mathcal{C}$  be a chain in  $(\mathcal{F}, \leq)$ , and define  $S^\infty = \{s \in \bigcup_{S \in \mathcal{C}} S \mid \text{there exists } S \in \mathcal{C} \text{ such that } s \in T \text{ whenever } T \in \mathcal{C} \text{ and } T \geq S\}$ . ( $S^\infty$  consists of all vertices of  $D$  that belong to every member of  $\mathcal{C}$  from some point on.)

We will prove that  $S^\infty$  is an upper bound of  $\mathcal{C}$ .

(2.2.1)  $S^\infty \neq \emptyset$ , and for each  $S \in \mathcal{C}$ ,  $S^\infty \geq S$ .

Let  $S \in \mathcal{C}$  and  $t_0 \in S$ . We will prove that there exists  $t \in S^\infty$  such that either  $t_0 = t$  or  $(t_0 \xrightarrow{1} t$  and  $t \nrightarrow t_0)$ . If  $t_0 \in S^\infty$  we are done. So assume  $t_0 \notin S^\infty$ . We proceed by contradiction; suppose that if  $t \in V(D)$  with  $(t_0 \xrightarrow{1} t$  and  $t \nrightarrow t_0)$ , then  $t \notin S^\infty$ . Take  $T_0 = S$ . Since  $t_0 \notin S^\infty$  we have that there exists  $T_1 \in \mathcal{C}$ ,  $T_1 \geq T_0$  such that  $t_0 \notin T_1$ . Hence there exists  $t_1 \in T_1$  such that  $t_0 \xrightarrow{1} t_1$  and  $t_1 \nrightarrow t_0$ . And our assumption implies  $t_1 \notin S^\infty$ . The fact  $t_1 \notin S^\infty$  implies  $t_1 \notin T_2$  for some  $T_2 \in \mathcal{C}$ ,  $T_2 \geq T_1$ . Hence there exists  $t_2 \in T_2$  such that  $t_1 \xrightarrow{1} t_2$  and  $t_2 \nrightarrow t_1$ . Since  $D_1$  is a quasi-transitive digraph, it follows from Lemma 2.1 on the sequence  $\tau_2 = (t_0, t_1, t_2)$  that  $\tau_2$  is an asymmetrical directed path of  $D$  contained in  $D_1$ ,  $(t_0 \xrightarrow{1} t_2$  and  $t_2 \nrightarrow t_0)$ ; and  $t_2 \notin S^\infty$ . We may continue this way and we obtain, for each  $n \in \mathbb{N}$ ,  $T_n \in \mathcal{C}$ ,  $t_n \in T_n$ ,  $(t_0 \xrightarrow{1} t_n$  and  $t_n \nrightarrow t_0)$  and  $t_n \notin S^\infty$ . Hence there exists  $T_{n+1} \in \mathcal{C}$  such that  $T_{n+1} \geq T_n$  and  $t_n \notin T_{n+1}$ . So there exists  $t_{n+1} \in T_{n+1}$  with  $(t_n \xrightarrow{1} t_{n+1}$  and  $t_{n+1} \nrightarrow t_n)$ .

Since  $D_1$  is a quasi-transitive digraph, and  $(t_n \xrightarrow{1} t_{n+1}$  and  $t_{n+1} \nrightarrow t_n)$  for each  $n \in \mathbb{N}$ , it follows from Lemma 2.1 (on the sequence  $\tau_{n+1} = (t_0, t_1, \dots, t_{n+1})$ ) that  $\tau_{n+1}$  is an asymmetrical directed path contained in  $D_1$  and in particular  $(t_0 \xrightarrow{1} t_{n+1}$  and  $t_{n+1} \nrightarrow t_0)$ . Our assumption implies  $t_{n+1} \notin S^\infty$ . Now consider the sequence  $\tau = (t_n)_{n \in \mathbb{N}}$ . For each  $n \in \mathbb{N}$  we have  $(t_n \xrightarrow{1} t_{n+1}$  and  $t_{n+1} \nrightarrow t_n)$ , and observe that for  $n < m$ ,  $\{t_n, t_m\} \subseteq V(\tau_m)$ , and since  $\tau_m$  is a directed path we have  $t_n \neq t_m$ . Hence  $\tau$  is an asymmetrical infinite outward path contained in  $D_1$ , a contradiction. We conclude that there exists  $t \in S^\infty$  such that  $(t_0 \xrightarrow{1} t$  and  $t \nrightarrow t_0)$ .

(2.2.2)  $S^\infty$  is an independent set.

Let  $s_1, s_2 \in S^\infty$  and suppose without loss of generality that  $S_1, S_2 \in \mathcal{C}$  are such that  $s_1 \in S_1, s_2 \in S_2, S_1 \leq S_2$ , since  $s_1 \in S^\infty$  we have  $s_1 \in S$  whenever  $S \in \mathcal{C}$  and  $S \geq S_1$ , so  $s_1 \in S_2$ , and since  $S_2$  is independent, there is no arc in  $D$  between  $s_1$  and  $s_2$ .

(2.2.3)  $S^\infty \in \mathcal{F}$ .

Suppose  $S^\infty \xrightarrow{2} y$  with  $y \in V(D)$ , so there exists  $s \in S^\infty$  with  $s \xrightarrow{2} y$ . Let  $S \in \mathcal{C}$  such that  $s \in T$  for all  $T \in \mathcal{C}$ ,  $T \geq S$ . Since  $S \in \mathcal{F}$  we have  $y \rightarrow S$ , so there exists  $s' \in S$  with  $y \rightarrow s'$ . When  $s' \in S^\infty$  we are done. When  $s' \notin S^\infty$  we analyze the two possibilities;  $y \xrightarrow{1} s'$  or  $y \xrightarrow{2} s'$ . First suppose  $y \xrightarrow{2} s'$ . Since  $s \xrightarrow{2} y$  and  $D_2$  is a quasi-transitive digraph it follows that  $s \xrightarrow{2} s'$  or  $s' \xrightarrow{2} s$  which is impossible as  $S$  is an independent set and  $\{s, s'\} \subseteq S$ . Now suppose  $y \xrightarrow{1} s'$ . Since  $s' \in S, S \leq S^\infty$  by (2.2.1) and  $s' \notin S^\infty$ , there exists  $t \in S^\infty$  such that  $s' \xrightarrow{1} t$  and  $t \nrightarrow s'$ . So we obtain  $y \xrightarrow{1} t$  or  $t \xrightarrow{1} y$  (as  $y \xrightarrow{1} s', s' \xrightarrow{1} t$  and  $D_1$  is a quasi-transitive digraph). If  $y \xrightarrow{1} t$  then  $y \xrightarrow{1} S^\infty$  and we are done. If  $t \xrightarrow{1} y$  then we obtain the triangle  $(y, s', t, y)$  and it follows from the hypothesis that it has two symmetrical arcs and since  $t \nrightarrow s'$  we have  $s' \rightarrow y$  and  $y \rightarrow t$ , so  $y \rightarrow S^\infty$ .

We have proven that any chain in  $\mathcal{F}$  has an upper bound in  $\mathcal{F}$ , and so by Zorn's Lemma,  $(\mathcal{F}, \leq)$  contains maximal elements. Let  $S$  be a maximal element of  $(\mathcal{F}, \leq)$ .

(3)  $S$  is a kernel of  $D$ .

Since  $S \in \mathcal{F}$ ,  $S$  is an independent set of vertices of  $D$ .

(3.1) For each  $x \in (V(D) - S)$  there exists an  $xS$ -arc.

Suppose by contradiction there exists  $x \in (V(D) - S)$  such that  $x \nrightarrow S$ .

**(3.1.1)** There exists a vertex  $x_0 \in V(D)$  such that  $x_0 \not\rightarrow S$  and  $x_0$  satisfies:  $x_0 \xrightarrow{2} y$  and  $y \not\rightarrow S$  imply  $y \rightarrow x_0$  for all vertices  $y \in V(D)$ . Let  $U = \{z \in V(D_2) - S \mid z \not\rightarrow S\}$ . When  $U \neq \emptyset$ , it follows from Lemma 2.2 (applied on  $D_2$  and  $U$ ) that there exists  $x_0$  with the required properties. When  $U = \emptyset$  it follows from our assumption that  $z \not\rightarrow S$ , for some vertex  $z$  in  $V(D_1) - (S \cup V(D_2))$ , and we take  $x_0$  to be any such vertex.

Note that the choice of  $x_0$  implies  $x_0 \not\rightarrow S$  and since  $S \in \mathcal{F}$ , we also have  $S \xrightarrow{2} x_0$ . Let  $T = \{s \in S \mid s \xrightarrow{1} x_0\}$ , it follows from above that  $T \cup \{x_0\}$  is an independent set of vertices of  $D$ .

**(3.1.2)**  $T \cup \{x_0\} \in \mathcal{F}$ .

Suppose  $T \cup \{x_0\} \xrightarrow{2} y$  and  $y \not\rightarrow T$ . We will prove  $y \rightarrow x_0$ . First we make the following observation.

**(3.1.2.1)** If  $y \xrightarrow{1} (S - T)$  then  $y \rightarrow x_0$ .

Let  $s \in (S - T)$  such that  $y \xrightarrow{1} s$ . Since  $s \in (S - T)$  we have  $s \xrightarrow{1} x_0$ . Now the fact that  $D_1$  is a quasi-transitive digraph implies  $y \xrightarrow{1} x_0$  or  $x_0 \xrightarrow{1} y$ . If  $x_0 \xrightarrow{1} y$  then  $(y, s, x_0, y)$  is a triangle which by the hypothesis has two symmetrical arcs, and since  $x_0 \not\rightarrow s$  it follows that  $y \rightarrow x_0$ .

We proceed to prove (3.1.2) by considering the two following cases:

*Case a:*  $T \xrightarrow{2} y$ .

Since  $T \subset S$  we have  $S \xrightarrow{2} y$  and the fact  $S \in \mathcal{F}$  implies  $y \rightarrow S$ . So  $y \rightarrow (S - T)$  (as we are assuming  $y \not\rightarrow T$ ).

When  $y \xrightarrow{1} (S - T)$  it follows from (3.1.2.1) that  $y \rightarrow x_0$ .

When  $y \xrightarrow{2} (S - T)$ , since we have  $T \xrightarrow{2} y$  and  $D_2$  is a quasi-transitive digraph, we obtain  $T \xrightarrow{2} (S - T)$  or  $(S - T) \xrightarrow{2} T$  and this is impossible as  $T \subseteq S$  and  $S$  is an independent set.

*Case b:*  $x_0 \xrightarrow{2} y$ .

We consider two possible subcases:

*Case b.1:*  $y \not\rightarrow S$ .

Since  $x_0 \xrightarrow{2} y$  and  $y \not\rightarrow S$ , the choice of  $x_0$  (see (3.1.1)) implies  $y \rightarrow x_0$ .

*Case b.2:*  $y \rightarrow S$ .

In this case we have  $y \rightarrow (S - T)$  (as we are assuming  $y \not\rightarrow T$ ).

When  $y \xrightarrow{2} (S - T)$ , since  $x_0 \xrightarrow{2} y$  and  $D_2$  is a quasi-transitive digraph, we have  $x_0 \xrightarrow{2} (S - T)$  or  $(S - T) \xrightarrow{2} x_0$ . Now recalling  $x_0 \not\rightarrow S$ , we obtain  $(S - T) \xrightarrow{2} x_0$  and since  $S \in \mathcal{F}$  it follows  $x_0 \rightarrow S$  which is impossible.

When  $y \xrightarrow{1} (S - T)$  it follows from (3.1.2.1) that  $y \rightarrow x_0$ .

**(3.1.3)**  $S < T \cup \{x_0\}$ .

For  $s \in (S - T)$  we have  $s \xrightarrow{1} x_0$  and we have noted  $x_0 \not\rightarrow S$ ; hence  $S \leq T \cup \{x_0\}$ . Moreover since  $x_0 \notin S$  (by construction in (3.1.1)) we have  $S < T \cup \{x_0\}$ .

Clearly propositions (3.1.2) and (3.1.3) contradict that  $S$  is a maximal element of  $(\mathcal{F}, \leq)$ .  $\square$

**Remark 2.4.** The condition that  $D_i$  has no infinite outward path in Theorem 2.3 is necessary.

Consider the following digraph  $D'$  with  $V(D') = \{u_n \mid n \in \mathbb{N}\}$  and  $A(D') = \{(u_n, u_m) \mid n, m \in \mathbb{N} \text{ and } n < m\}$ ,  $D_1 = D'$ ,  $D_2 = D'$  and  $D = D_1 \cup D_2$ .

**Remark 2.5.** The following digraph  $D$  is the union of two quasi-transitive finite digraphs; each triangle in  $D$  has at least one symmetrical arc and  $D$  has no kernel.

$$V(D_1) = \{u_0, u_1, u_2, u_3\},$$

$$V(D_2) = V(D_1) \cup \{w\},$$

$$A(D_1) = \{(u_i, u_{i+1}) \mid i \in \{0, 1, 2, 3\} \pmod{4}\} \cup \{(u_0, u_2), (u_2, u_0), (u_1, u_3), (u_3, u_0)\},$$

$$A(D_2) = \{(w, u_i) \mid i \in \{0, 1, 2, 3\}\},$$

$$D = D_1 \cup D_2.$$

**Remark 2.6.** Clearly  $\vec{C}_5$  the directed cycle of length 5 is the union of two finite digraphs,  $\vec{C}_5$  has no triangle and  $\vec{C}_5$  has no kernel.

We conclude that the conditions on Theorem 2.3 are tight.

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